

Generics und Wildcards in Java 5.0

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Overview

Motivation

Simple types

Subtyping

Conclusion

Type terms in Java 5.0:

explicitly used:

```
Vector<Vector<Integer>>  
Vector<? extends List<Object>>  
Vector<? super List<Object>>
```

only inferred:

```
? extends List<Object>  
? super List<Object>
```

Formal description of type terms

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Definition *Rank alphabet:*

- ▶ $\Theta = \Theta_{n \in \mathbb{N}}^{(n)}$ (rank alphabet)
- ▶ TV (set of type variables)
- ▶ $T_{\Theta}(TV)$ (Java 5.0 type terms)

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Example:

```
class A<a> implements I<a> { ... }
class B<a> extends A<a> { ... }
class C<a extends I<b>,b> { ... }
interface I<a> { ... }
```

- ▶ $\Theta^{(1)} = \{A, B, I\}$
 $\Theta^{(2)} = \{C\}$
- ▶ $A<Integer>, A<B<Boolean>>, C<A<Object>, Object> \in T_{\Theta}(TV)$

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- ▶ $\Theta^{(1)} = \{A, B, I\}$
 $\Theta^{(2)} = \{C\}$
- ▶ $A<Integer>, A<B<Boolean>>, C<A<Object>, Object> \in T_{\Theta}(TV)$
- ▶ $C<C<a>, a> \in T_{\Theta}(TV)$, **but no type term!!!**

Bounded type variables

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Example:

```
class BoundedTypeVars<a extends Number> {
    <t extends Vector<Integer> & J<a> & I,
    r extends Number> void m ( ... ) { ... } }
```

- ▶ $BTV^{(Number)} = \{ a, r \}$
- ▶ $BTV^{(Vector<Integer> & J<a> & I)} = \{ t \}$

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Definition *Type signature*:

- ▶ $TS = (TC^{(tv_1 \dots tv_n)})_{tv_i \in BTV}$ (BTV^* -indexed set of *type constructors*)

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Example:

```
class A<a> implements I<a> { ... }
class B<a> extends A<a> { ... }
class C<a extends I<b>, b> { ... }
interface I<a> { ... }
interface J<a> { ... }
class D<a extends B<a> & J<b>, b> { ... }
```

- ▶ $TC^{(a|_{\text{Object}})} = \{A, B, I, J\}$
- ▶ $TC^{(a|_{I} \ b|_{\text{Object}})} = \{C\}$
- ▶ $TC^{(a|_{B<a> \& \ J} \ b|_{\text{Object}})} = \{D\}$

Simple types $S\text{Type}_{TS}(BTV)$ (first approach)

- ▶ $BTV \subseteq S\text{Type}_{TS}(BTV)$
- ▶ $TC^() \subseteq S\text{Type}_{TS}(BTV)$
- ▶ For $C \in TC^{(a_1|ty_1 \dots a_n|ty_n)}$ and a substitution σ

$$C\langle\sigma(a_1), \dots, \sigma(a_n)\rangle \in S\text{Type}_{TS}(BTV),$$

if $\sigma(a_i) \leq^* \sigma(ty_i)$ (\leq^* is the subtyping relation).

Example:

```
class A<a> implements I<a> { ... }
```

```
class B<a> extends A<a> { ... }
```

```
class C<a extends I<b>, b> { ... }
```

```
interface I<a> { ... }
```

- ▶ $TC(a|_{\text{Object}}) = \{A, B, I\}$
 $TC(a|_{I} b|_{\text{Object}}) = \{C\}$
- ▶ $A<Integer>, A<B<Boolean>>, C<A<Object>, Object>$
 $\in SType_{TS}(BTV)$
- ▶ $C<C<a>, a> \notin SType_{TS}(BTV)$

Subtyping ordering \leq^*

Definition:

- ▶ if θ extends/implements θ' then $\theta \leq^* \theta'$.
- ▶ if $\theta_1 \leq^* \theta_2$ then $\sigma_1(\theta_1) \leq^* \sigma_2(\theta_2)$ for all substitutions σ_1, σ_2 with $\sigma_1(a) = \sigma_2(a)$ (soundness condition).

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class B<a> extends A<a> { ... }
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interface I<a> { ... }
```

- ▶ $A<a> \leq^* I<a>$, as $A<a>$ implements $I<a>$
- ▶ $A<Integer> \leq^* I<Integer>$, where $\sigma_1 = [a \mapsto Integer] = \sigma_2$
- ▶ $A<Integer> \not\leq^* I<Object>$, although indeed $Integer \leq^* Object$, but $\sigma_1(a) \neq \sigma_2(a)$.

Soundness condition: $\sigma_1(a) = \sigma_2(a)$

```

class Super { ... }
class Sub extends Super { ... }

class Application {
  public static void main(String[] args) {
    Vector<Super> v = new Vector<Sub> (); //not allowed
    v.addElement(new Super());        //wrong: Super  $\not\leq^*$  Sub
  }
}

```

Soundness condition: $\sigma_1(a) = \sigma_2(a)$

```

class Super { ... }
class Sub extends Super { ... }

class Application {
  public static void main(String[] args) {
    Vector<Super> v = new Vector<Sub> (); //not allowed
    v.addElement(new Super()); //wrong: Super  $\not\leq^*$  Sub
  }

  void meth (Vector<Sub> subvec) {
    Vector<Super> supervec = subvec; //could make sense
    ... //but not allowed
  }
}

```

Introduction of wildcards

extends-wildcards

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Means: The elements of the vector `v` are `subtypes` of `Super`.

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- ▶ $\text{Vector}\langle\text{Sub}\rangle \leq^* \text{Vector}\langle? \text{ extends Super}\rangle$
- ▶ $\text{Super} \not\leq^* ? \text{ extends Super}$

Introduction of wildcards

extends-wildcards

```
Vector<? extends Super> v = new Vector<Sub> ();
```

Means: The elements of the vector `v` are `subtypes` of `Super`.

⇒

- ▶ `Vector<Sub> ≤* Vector<? extends Super>`
- ▶ `Super ≤* ? extends Super`
- ▶ `v.addElement(new Super());` **not valid**

super-wildcards

Problem:

```
void addSuperElement(Vector<Sub> v) {  
    v.addElement(new Sub ());  
}
```

Application of `addSuperElement(new Vector<Super> ());` is not allowed, although it make sense.

super-wildcards

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```
void addSuperElement(Vector<Sub> v) {  
    v.addElement(new Sub ());  
}
```

Application of `addSuperElement(new Vector<Super> ());` is not allowed, although it make sense.

Solution:

```
void addSuperElement(Vector<? super Sub> v) {  
    v.addElement(new Sub ());  
}
```

Now `addSuperElement(new Vector<Super> ());` is allowed.

Properties of super-wildcards

- ▶ $\text{Vector}\langle\text{Super}\rangle \leq^* \text{Vector}\langle? \text{ super Sub}\rangle$, as $\text{Sub} \leq^* \text{Super}$.
- ▶ $? \text{ super Sub} \not\leq^* \text{Sub}$
- ▶ `Vector<? super Sub> v;`
`Sub e = v.elementAt(0);` **not valid.**

Abbreviations for wildcards:

Instead of `A<? extends B>` we write

`A<?B>`

and instead of `C<? super D>` we write

`C<?D>`.

Nested extends–declarations

```
class Matrix<a> extends Vector<Vector<a>> { }
```

From $\text{Matrix}\langle a \rangle \leq^* \text{Vector}\langle \text{Vector}\langle a \rangle \rangle$ follows

▶ $\text{Matrix}\langle \text{Super} \rangle \leq^* \text{Vector}\langle \text{Vector}\langle \text{Super} \rangle \rangle$

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Is the following correct?

- ▶ $\text{Matrix}\langle \text{Sub} \rangle \leq^* \text{Vector}\langle \text{Vector}\langle ? \text{Super} \rangle \rangle$

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Vector<Vector<? extends Super>> v = new Matrix<Sub>();
v.addElement(new Vector<Super>());
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Vector<Vector<? extends Super>> v = new Matrix<Sub>();
v.addElement(new Vector<Super>());
//wrong: Vector<Super>  $\not\leq^*$  Vector<Sub>
```

⇒ $\text{Matrix}\langle \text{Sub} \rangle \not\leq^* \text{Vector}\langle \text{Vector}\langle ?\text{Super} \rangle \rangle$

Capture conversion (for type constructors without bounds)

Let $C \in TC^{(a_1|_{\text{Object}}, \dots, a_n|_{\text{Object}})}$.

$CC(C\langle\theta_1, \dots, \theta_n\rangle) = C\langle\bar{\theta}_1, \dots, \bar{\theta}_n\rangle$ is defined as:

- ▶ if $\theta_i = ?\theta'_i$: $\bar{\theta}_i = b_i|\theta'_i$, b_i fresh type variable with *upper bound* θ'_i .
- ▶ if $\theta_i = \theta'_i?$: $\bar{\theta}_i = \theta'_i|b_i$, b_i fresh type variable with *lower bound* θ'_i .
- ▶ otherwise $\bar{\theta}_i = \theta_i$

Capture conversion (for type constructors with bounds)

Let $C \in TC(a_1|u_1, \dots, a_n|u_n)$.

$CC(C\langle\theta_1, \dots, \theta_n\rangle) = C\langle\bar{\theta}_1, \dots, \bar{\theta}_n\rangle$ is defined as:

- ▶ if $\theta_i = ?$: $\bar{\theta}_i = b_i|u_i[a_j \mapsto \bar{\theta}_j \mid 1 \leq j \leq n]$, b_i fresh type variable
- ▶ if $\theta_i = ?\theta'_i$: $\bar{\theta}_i = b_i|\theta'_i \& u_i[a_j \mapsto \bar{\theta}_j \mid 1 \leq j \leq n]$, b_i fresh type variable with *upper bound* $\theta'_i \& u_i[a_j \mapsto \bar{\theta}_j \mid 1 \leq j \leq n]$.
- ▶ if $\theta_i = \theta'_i$: $\bar{\theta}_i = \theta'_i|b_i|u_i[a_j \mapsto \bar{\theta}_j \mid 1 \leq j \leq n]$, b_i fresh type variable with *lower bound* θ'_i and *upper bound* $u_i[a_j \mapsto \bar{\theta}_j \mid 1 \leq j \leq n]$.
- ▶ otherwise $\bar{\theta}_i = \theta_i$

Simple types $S\text{Type}_{TS}(BTV)$ (t.c. without bounds)

- ▶ $BTV \subseteq S\text{Type}_{TS}(BTV)$
 - ▶ $TC^{()}\subseteq S\text{Type}_{TS}(BTV)$
 - ▶ For $ty_i \in S\text{Type}_{TS}(BTV)$
 - $\cup \{?\}$
 - $\cup \{? \text{ extends } \tau \mid \tau \in S\text{Type}_{TS}(BTV)\}$
 - $\cup \{? \text{ super } \tau \mid \tau \in S\text{Type}_{TS}(BTV)\}$
- and $C \in TC^{(a_1|\text{Object}\dots a_n|\text{Object})}$ holds

$$C\langle ty_1, \dots, ty_n \rangle \in S\text{Type}_{TS}(BTV).$$

Simple types $\text{SType}_{TS}(BTV)$ (t.c. with bounds)

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and $C \in TC^{(a_1|b_1 \dots a_n|b_n)}$ holds

$$C \langle ty_1, \dots, ty_n \rangle \in \text{SType}_{TS}(BTV)$$

if for $CC(C \langle ty_1, \dots, ty_n \rangle) = C \langle \overline{ty_1}, \dots, \overline{ty_n} \rangle$ holds:

$$\overline{ty_i} \leq^* b_i [a_j \mapsto \overline{ty_j} \mid 1 \leq j \leq n],$$

Subtyping ordering \leq^* (extension)

- ▶ if θ extends/implements θ' then $\theta \leq^* \theta'$.
- ▶ if $\theta_1 \leq^* \theta_2$ then $\sigma_1(\theta_1) \leq^* \sigma_2(\theta_2)$, where for each type variable a of θ_2 holds $\sigma_1(a) = \sigma_2(a) \in \text{SType}_{TS}(\text{BTV})$ (soundness condition).
- ▶ It holds $C\langle\theta_1, \dots, \theta_n\rangle \leq^* C\langle\theta'_1, \dots, \theta'_n\rangle$ if for θ_i and θ'_i either
 - ▶ $\theta_i = ?\bar{\theta}_i$, $\theta'_i = ?\bar{\theta}'_i$ and $\bar{\theta}_i \leq^* \bar{\theta}'_i$ or
 - ▶ $\theta_i = ?\bar{\theta}_i$, $\theta'_i = ?\bar{\theta}'_i$ and $\bar{\theta}'_i \leq^* \bar{\theta}_i$ or
 - ▶ θ_i, θ'_i are no wildcard arguments and $\theta_i = \theta'_i$ or
 - ▶ $\theta'_i = ?\theta_i$ or
 - ▶ $\theta'_i = ?\theta_i$
- ▶ From $C\langle\bar{\theta}_1, \dots, \bar{\theta}_n\rangle \leq^* C\langle\theta'_1, \dots, \theta'_n\rangle$ follows with $C\langle\bar{\theta}_1, \dots, \bar{\theta}_n\rangle = CC(C\langle\theta_1, \dots, \theta_n\rangle)$: $C\langle\theta_1, \dots, \theta_n\rangle \leq^* C\langle\theta'_1, \dots, \theta'_n\rangle$
- ▶ $T|(\theta_1 \& \dots \& \theta_n) \leq^* \theta_i$ for any $1 \leq i \leq n$.
- ▶ $\theta \leq^* \theta | T$

Example: $\text{Matrix}\langle a \rangle \leq^* \text{Vector}\langle \text{Vector}\langle a \rangle \rangle$

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`Vector<Vector<Integer>>`



`Matrix<Integer>`

Example: $\text{Matrix}\langle a \rangle \leq^* \text{Vector}\langle \text{Vector}\langle a \rangle \rangle$

`Vector<Vector<Integer>>`

`Matrix<Integer>`

`Vector<?Vector<?Integer>>`

`Vector<?Vector<?X|Integer>>`

`Vector<?Vector<X|Integer>>`

`Vector<Vector<X|Integer>>`

`Matrix<X|Integer>`

`Matrix<?Integer>`

cc

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`Matrix<X|Integer>`

`Matrix<?Integer>`

cc

Motivation for a continuation of \leq^* on wildcard types

An element should be read from a vector of a wildcard type:

```
Vector<? extends Super> v = new Vector<Sub> ();
Super superElement = v.elementAt(i);
```

$\implies ? \text{ extends Super} \leq^* \text{Super}$

Motivation for a continuation of \leq^* on wildcard types

An element should be read from a vector of a wildcard type:

```
Vector<? extends Super> v = new Vector<Sub> ();
Super superElement = v.elementAt(i);
```

$\implies ? \text{ extends Super} \leq^* \text{ Super}$

An element of a **subclass** should be added to a vector of a **superclass**:

```
Vector<? super Super> v2 = new Vector<Super> ();
v2.addElement(new Sub());
```

$\implies \text{ Super, Sub} \leq^* ? \text{ super Super}$

Extended simple types

$\text{SType}_{TS}(BTV)$ is a set of simple types:

The corresponding set of *extended simple types* is given as

$$\begin{aligned} \text{ExtSType}_{TS}(BTV) = & \text{SType}_{TS}(BTV) \quad . \\ & \cup \{?\} \\ & \cup \{? \text{ extends } \theta \mid \theta \in \text{SType}_{TS}(BTV)\} \\ & \cup \{? \text{ super } \theta \mid \theta \in \text{SType}_{TS}(BTV)\} \end{aligned}$$

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$$\begin{aligned} \text{ExtSType}_{TS}(BTV) = & \text{SType}_{TS}(BTV) \\ & \cup \{?\} \\ & \cup \{? \text{ extends } \theta \mid \theta \in \text{SType}_{TS}(BTV)\} \\ & \cup \{? \text{ super } \theta \mid \theta \in \text{SType}_{TS}(BTV)\} \end{aligned} .$$

Wildcard types cannot be used, explicitly. They are only inferred.

\leq^* on $\text{ExtSType}_{TS}(BTV)$

For $\theta, \theta' \in \text{SType}_{TS}(BTV)$ with $\theta \leq^* \theta'$ holds:

- ▶ $\theta \leq^* ?\theta'$,
- ▶ $? \theta \leq^* \theta'$, and
- ▶ $? \theta \leq^* ? \theta'$.

Properties of type constructor applying

For $\theta \leq^* \theta'$:

- ▶ $C1\langle\theta\rangle \not\leq^* C1\langle\theta'\rangle$
- ▶ $C1\langle\theta'\rangle \not\leq^* C1\langle\theta\rangle$

Properties of type constructor applying

For $\theta \leq^* \theta'$:

- ▶ $C1\langle\theta\rangle \not\leq^* C1\langle\theta'\rangle$
- ▶ $C1\langle\theta'\rangle \not\leq^* C1\langle\theta\rangle$
- ▶ It holds

? extends $\theta \leq^* \theta'$

but

$C1\langle\theta\rangle \leq^* C1\langle ? \text{ extends } \theta' \rangle$

Properties of type constructor applying

For $\theta \leq^* \theta'$:

- ▶ $CL\langle\theta\rangle \not\leq^* CL\langle\theta'\rangle$
- ▶ $CL\langle\theta'\rangle \not\leq^* CL\langle\theta\rangle$
- ▶ It holds

? extends $\theta \leq^* \theta'$

but

$CL\langle\theta\rangle \leq^* CL\langle ? \text{ extends } \theta' \rangle$

- ▶ It holds

(? extends) $\theta \leq^* ? \text{ super } \theta'$

but

$CL\langle\theta'\rangle \leq^* CL\langle ? \text{ super } \theta \rangle$

Conclusion

- ▶ Formalization of the Java 5.0 type system
- ▶ Subtyping ordering on $S\text{Type}_{TS}(BTV)$ and $\text{ExtS}\text{Type}_{TS}(BTV)$
- ▶ Base for the type inference algorithm, which works, if
 - ▶ there is a subtyping ordering and
 - ▶ there is a type unification algorithm